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SHUMAKER & SIEFFERT, P.A. 1625 RADIO DRIVE, SUITE 300 WOODBURY, MN 55125			EXAMINER RICHARDSON, THOMAS W	
			ART UNIT 2444	PAPER NUMBER
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**Please find below and/or attached an Office communication concerning this application or proceeding.**

The time period for reply, if any, is set in the attached communication.

Notice of the Office communication was sent electronically on above-indicated "Notification Date" to the following e-mail address(es):

pairedocketing@ssiplaw.com

### Office Action Summary

**Application No.**

10/687,989

**Applicant(s)**

PILLAY ESNAULT, PADMA

**Examiner**

THOMAS RICHARDSON

**Art Unit**

2444

**Period for Reply** -- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE 3 MONTH(S) OR THIRTY (30) DAYS, WHICHEVER IS LONGER, FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133). Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

**Status**

- 1) ☒ Responsive to communication(s) filed on 06 May 2010.
- 2a) ☒ This action is **FINAL**. 2b) ☐ This action is non-final.
- 3) ☐ Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

**Disposition of Claims**

- 4) ☒ Claim(s) 1,6,8,11-13,15,18-22,26,27,30,32,33,39 and 40 is/are pending in the application.
- 4a) Of the above claim(s) \_\_\_\_\_ is/are withdrawn from consideration.
- 5) ☐ Claim(s) \_\_\_\_\_ is/are allowed.
- 6) ☒ Claim(s) 1,6,8,11-13,15,18-22,26,27,30,32,33,39 and 40 is/are rejected.
- 7) ☐ Claim(s) \_\_\_\_\_ is/are objected to.
- 8) ☐ Claim(s) \_\_\_\_\_ are subject to restriction and/or election requirement.

**Application Papers**

- 9) ☐ The specification is objected to by the Examiner.
- 10) ☐ The drawing(s) filed on \_\_\_\_\_ is/are: a) ☐ accepted or b) ☐ objected to by the Examiner.  
Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).  
Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).
- 11) ☐ The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

**Priority under 35 U.S.C. § 119**

- 12) ☐ Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).
- a) ☐ All b) ☐ Some \* c) ☐ None of:
1. ☐ Certified copies of the priority documents have been received.
  2. ☐ Certified copies of the priority documents have been received in Application No. \_\_\_\_\_.
  3. ☐ Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).

\* See the attached detailed Office action for a list of the certified copies not received.

**Attachment(s)**

- 1) ☐ Notice of References Cited (PTO-892)
- 2) ☐ Notice of Draftsperson's Patent Drawing Review (PTO-940)
- 3) ☐ Information Disclosure Statement(s) (PTO/SB/08)  
Paper No(s)/Mail Date \_\_\_\_\_
- 4) ☐ Interview Summary (PTO-413)  
Paper No(s)/Mail Date \_\_\_\_\_
- 5) ☐ Notice of Informal Patent Application
- 6) ☐ Other: \_\_\_\_\_

### **DETAILED ACTION**

Claims 1, 6, 8, 11-13, 15, 18-22, 26, 27, 30, 32, 33, 39, and 40 are pending for examination.

Claim 8 is amended.

Claims 1, 6, 8, 11-13, 15, 18-22, 26, 27, 30, 32, 33, 39, and 40 are rejected.

### ***Response to Arguments***

1. Applicant's arguments filed with respect to claim 11 have been fully considered but they are not persuasive. Applicant argues that cited references Pesce (US 7 328 278), Sankaran (US 2003/0231587), and Gaddis (US 7 554 930) do not teach the limitations of the claim. Examiner disagrees.
2. As per claim 11, applicant argues that cited references do not teach utilizing a metric. Both Gaddis and Sankaran teach utilizing a path metric such as AS\_Path (Gaddis column 6, lines 44-58, Sankaran paragraph 32), in addition to the metric utilized by Pesce to determine table information and which routes to maintain within the table. As such, it would have been obvious to one of ordinary skill in the art at the time of the invention to utilize metrics for links such as taught by Gaddis and Sankaran as the metric described by Pesce, as AS\_Path is a metric known to be used in the art of network routing (Gaddis, column 6, lines 44-46). Applicant further argues that cited references do not teach clearing routes from the table when a limit is reached. As cited, Sankaran teaches removing redundant routes from the table, thus clearing those routes from the table, and as a result, the table being altered from its original form as a result

of the threshold being reached (paragraph 35). As such, the previous rejection of claim 11 is maintained.

3. With respect to claims 1, 8, 18, 27, and 33, applicant relies on arguments similar to those presented with respect to claim 11, and as such, examiner relies on arguments presented in response. The previous rejections of claims 1, 8, 18, 27, and 33 are maintained.

***Claim Rejections - 35 USC § 103***

4. The text of those sections of Title 35, U.S. Code not included in this action can be found in a prior Office action.

5. Claims 1, 8, 11-13, 15, 18-22, 26, 27, 30, 33, 39, and 40 are rejected under 35 U.S.C. 103(a) as being unpatentable over US 7 328 278, Pesce et al, US 2003/0231587, Sankaran et al, and US 7 554 930, Gaddis et al.

6. As per claim 1, Pesce teaches a method comprising:

routes exported from an exterior routing protocol executing on a network device to an interior routing protocol executing on the network device (column 2, lines 51-64, where the network element may contain multiple routing databases including those that may run both interior and exterior routing protocols, also column 3, lines 23-28, where the routing information may be exported from one database operating in one protocol to another database operating in another protocol);

maintaining routes exported from the exterior routing protocol executing on the network device to the interior routing protocol executing on the network device (column 2, lines 51-64, where the network element may contain multiple routing databases

including those that may run both interior and exterior routing protocols, also column 3, lines 23-28, where the routing information may be exported from one database operating in one protocol to another database operating in another protocol); and

blocking routes exported from the exterior routing protocol to the interior routing protocol (column 4, lines 52-61, where the updated mappings may be blocked by the manager);

updating routing information to associate the routes with a maximum metric that defines a maximum distance from the network device to neighboring network devices (column 5, lines 14-16, where each entry in the table may be associated with a metric value and scale, also column 11, lines 52-60, where routes may be associated with a metric); and

advertising the updated routing information to a network device (column 15, lines 49-51, where routes may be advertised).

Pesce does not expressly teach a client setting an export limit for the device. Sankaran teaches a method for managing discard algorithms comprising:

maintaining a count of routes exported (paragraph 35, where a threshold may be reached); and

rejecting additional routes exported when the count exceeds the export limit set by the command (paragraph 45, where additional routes may be discarded once a threshold is reached).

It would have been obvious to one of ordinary skill in the art at the time of the invention to utilize a route count and rejecting such as taught by Sankaran in a routing system

such as taught by Pesce. Pesce's system generally allows for routes to be exported from one protocol operating on a network element to another protocol operating on a network element. Pesce further allows for blocking routes by a manager (column 4, lines 52-61). It would be beneficial in such a system to utilize a threshold and count of routes for blocking, as the routing manager may operate under blocking conditions after a threshold of a table is reached, such as a percentage of the memory being used (paragraph 45).

Neither Pesce nor Sankaran teaches a user defined size threshold. Gaddis teaches a method for improving databases comprising:

receiving a command from a client to specify a limit for a table size (column 17, lines 29-34, where a prefix limit may be set for injected routes).

It would have been obvious to utilize a user defined limit such as that taught by Gaddis in a system for exporting routes such as that taught by either Pesce or Sankaran.

Pesce's system generally allows for routes to be exported according to definitions by a user, including those determining whether routes may be exported (column 7, lines 11-14). Pesce's system also allows a manager to determine when blocking conditions occur (column 4, lines 52-61). It would be beneficial to allow for a user to determine a maximum number of routes, such as taught by Sankaran, as the routing manager may operate under blocking conditions after a threshold of a table is reached, such as a percentage of the memory being used (paragraph 45). It would further be beneficial for a user to determine the maximum, such as taught by Gaddis, this allows the user to set further rules regarding route exporting, in addition to those as taught by Pesce.

7. As per claim 27, Pesce teaches a computer-readable medium comprising instructions to cause a processor to:

present a management interface to receive a command from a client (column 7, lines 11-14, where a user may set rules governing route exportation);

receive a command, from the client through the management interface, to specify rules for routes exported from an exterior routing protocol executing on a network device to an interior routing protocol executing on the network device (column 7, lines 11-14, where a user may set rules governing route exportation);

maintain rules of routes exported from an exterior routing protocol executing on the network device to a plurality of instances of an interior routing protocol executing on the network device (column 7, lines 11-14, where a user may set rules governing route exportation, also column 2, lines 51-64, where the network element may contain multiple routing databases including those that may run both interior and exterior routing protocols, also column 3, lines 23-28, where the routing information may be exported from one database operating in one protocol to another database operating in another protocol); and

identify one of the instances of the interior routing protocol to which the routes were exported (column 2, lines 51-64, where the network element may contain multiple routing databases including those that may run both interior and exterior routing protocols, also column 3, lines 23-28, where the routing information may be exported from one database operating in one protocol to another database operating in another protocol);

compare the rule for the identified one of the instances to an export limit (column 7, lines 11-14, where a user may set rules governing route exportation, also column 2, lines 51-64, where the network element may contain multiple routing databases including those that may run both interior and exterior routing protocols, also column 3, lines 23-28, where the routing information may be exported from one database operating in one protocol to another database operating in another protocol); and

reject additional routes to be exported from the exterior routing protocol to the identified one of the instances based on the comparison (column 4, lines 52-61, where the updated mappings may be blocked by the manager).

Pesce does not expressly teach a client setting an export limit for the device. Sankaran teaches a method for managing discard algorithms comprising:

maintaining a count of routes exported (paragraph 35, where a threshold may be reached); and

rejecting additional routes exported when the count exceeds the export limit set by the command (paragraph 45, where additional routes may be discarded once a threshold is reached).

It would have been obvious to one of ordinary skill in the art at the time of the invention to utilize a route count and rejecting such as taught by Sankaran in a routing system such as taught by Pesce. Pesce's system generally allows for routes to be exported from one protocol operating on a network element to another protocol operating on a network element. Pesce further allows for blocking routes by a manager (column 4, lines 52-61). It would be beneficial in such a system to utilize a threshold and count of



routes for blocking, as the routing manager may operate under blocking conditions after a threshold of a table is reached, such as a percentage of the memory being used (paragraph 45).

Neither Pesce nor Sankaran teaches a user defined size threshold. Gaddis teaches a method for improving databases comprising:

receiving a command from a client to specify a limit for a table size (column 17, lines 29-34, where a prefix limit may be set for injected routes).

It would have been obvious to utilize a user defined limit such as that taught by Gaddis in a system for exporting routes such as that taught by either Pesce or Sankaran.

Pesce's system generally allows for routes to be exported according to definitions by a user, including those determining whether routes may be exported (column 7, lines 11-14). Pesce's system also allows a manager to determine when blocking conditions occur (column 4, lines 52-61). It would be beneficial to allow for a user to determine a maximum number of routes, such as taught by Sankaran, as the routing manager may operate under blocking conditions after a threshold of a table is reached, such as a percentage of the memory being used (paragraph 45). It would further be beneficial for a user to determine the maximum, such as taught by Gaddis, this allows the user to set further rules regarding route exporting, in addition to those as taught by Pesce.

8. As per claim 30, Sankaran further teaches updating muting information to associate the routes with a maximum metric that defines a maximum distance from the network device to neighboring network devices when the count exceeds the export limit (paragraph 5, where the redundant routes may be preferenced by distance); and

advertising the updated routing information to a network device (paragraph 3, where the network may redefine paths using BGP or OSPF);

updating routing information when the count exceeds the limit to clear the routes exported (paragraph 45, where routes may be deleted from the table when a limit is reached or exceeded).

Pesce further teaches waiting for user intervention before accepting routes (column 7, lines 11-14, where a user may set rules governing route exportation).

9. As per claim 8, Pesca teaches a method comprising:

an export rule and an associated one of a plurality of instances of an interior routing protocol executing on a network device (column 7, lines 11-14, where a user may set rules governing route exportation, also column 2, lines 51-64, where the network element may contain multiple routing databases including those that may run both interior and exterior routing protocols, also column 3, lines 23-28, where the routing information may be exported from one database operating in one protocol to another database operating in another protocol);

maintaining rules of routes exported from an exterior routing protocol executing on the network device to each of the instances of the interior routing protocol executing on the network device (column 7, lines 11-14, where a user may set rules governing route exportation, also column 2, lines 51-64, where the network element may contain multiple routing databases including those that may run both interior and exterior routing protocols, also column 3, lines 23-28, where the routing information may be exported

from one database operating in one protocol to another database operating in another protocol);

identifying one of the instances of the interior routing protocol to which the routes were exported (column 7, lines 11-14, where a user may set rules governing route exportation, also column 2, lines 51-64, where the network element may contain multiple routing databases including those that may run both interior and exterior routing protocols, also column 3, lines 23-28, where the routing information may be exported from one database operating in one protocol to another database operating in another protocol);

comparing the rules for the identified one of the instances to the rule specified by the command for the identified instance (column 7, lines 11-14, where a user may set rules governing route exportation, also column 2, lines 51-64, where the network element may contain multiple routing databases including those that may run both interior and exterior routing protocols, also column 3, lines 23-28, where the routing information may be exported from one database operating in one protocol to another database operating in another protocol); and

rejecting additional routes exported to the identified one of the instances of the interior routing protocol based on the comparison when the rule determines (column 4, lines 52-61, where the updated mappings may be blocked by the manager).  
Pesce does not expressly teach a client setting an export limit for the device. Sankaran teaches a method for managing discard algorithms comprising:

maintaining a count of routes exported (paragraph 35, where a threshold may be reached); and

rejecting additional routes exported when the count exceeds the export limit set by the command (paragraph 45, where additional routes may be discarded once a threshold is reached).

It would have been obvious to one of ordinary skill in the art at the time of the invention to utilize a route count and rejecting such as taught by Sankaran in a routing system such as taught by Pesce. Pesce's system generally allows for routes to be exported from one protocol operating on a network element to another protocol operating on a network element. Pesce further allows for blocking routes by a manager (column 4, lines 52-61). It would be beneficial in such a system to utilize a threshold and count of routes for blocking, as the routing manager may operate under blocking conditions after a threshold of a table is reached, such as a percentage of the memory being used (paragraph 45).

Neither Pesce nor Sankaran teaches a user defined size threshold. Gaddis teaches a method for improving databases comprising:

receiving a prefix limit command from a client to specify a limit for a table size (column 17, lines 29-34, where a prefix limit may be set for injected routes).

It would have been obvious to utilize a user defined limit such as that taught by Gaddis in a system for exporting routes such as that taught by either Pesce or Sankaran. Pesce's system generally allows for routes to be exported according to definitions by a user, including those determining whether routes may be exported (column 7, lines 11-

14). Pesce's system also allows a manager to determine when blocking conditions occur (column 4, lines 52-61). It would be beneficial to allow for a user to determine a maximum number of routes, such as taught by Sankaran, as the routing manager may operate under blocking conditions after a threshold of a table is reached, such as a percentage of the memory being used (paragraph 45). It would further be beneficial for a user to determine the maximum, such as taught by Gaddis, this allows the user to set further rules regarding route exporting, in addition to those as taught by Pesce.

10. As per claim 11, Pesce teaches a method comprising:

exporting routes from an exterior routing protocol process executing on the network device to an interior routing protocol process executing on the network device (column 2, lines 51-64, where the network element may contain multiple routing databases including those that may run both interior and exterior routing protocols, also column 3, lines 23-28, where the routing information may be exported from one database operating in one protocol to another database operating in another protocol); and

the network device: (i) updates routing information of the interior routing protocol to clear the routes previously exported from the exterior routing protocol, (ii) rebuilds the routing information of the interior routing protocol by updating the routing information of the interior routing protocol to associate interior routes with a maximum metric that defines a maximum distance from the network device to neighboring network devices, and (iii) advertises the updated routing information to another network device (column 5,

lines 14-16, where each entry in the table may be associated with a metric value and scale, also column 11, lines 52-60, where routes may be associated with a metric).

Pesce does not expressly teach a client setting an export limit for the device. Sankaran teaches a method for managing discard algorithms comprising:

maintaining a count of routes exported (paragraph 35, where a threshold may be reached); and

rejecting additional routes exported when the count exceeds the export limit set by the command (paragraph 45, where additional routes may be discarded once a threshold is reached).

It would have been obvious to one of ordinary skill in the art at the time of the invention to utilize a route count and rejecting such as taught by Sankaran in a routing system such as taught by Pesce. Pesce's system generally allows for routes to be exported from one protocol operating on a network element to another protocol operating on a network element. Pesce further allows for blocking routes by a manager (column 4, lines 52-61). It would be beneficial in such a system to utilize a threshold and count of routes for blocking, as the routing manager may operate under blocking conditions after a threshold of a table is reached, such as a percentage of the memory being used (paragraph 45).

Neither Pesce nor Sankaran teaches a user defined size threshold. Gaddis teaches a method for improving databases comprising:

receiving a prefix limit command from a client to specify a limit for a table size (column 17, lines 29-34, where a prefix limit may be set for injected routes).

It would have been obvious to utilize a user defined limit such as that taught by Gaddis in a system for exporting routes such as that taught by either Pesce or Sankaran.

Pesce's system generally allows for routes to be exported according to definitions by a user, including those determining whether routes may be exported (column 7, lines 11-14). Pesce's system also allows a manager to determine when blocking conditions occur (column 4, lines 52-61). It would be beneficial to allow for a user to determine a maximum number of routes, such as taught by Sankaran, as the routing manager may operate under blocking conditions after a threshold of a table is reached, such as a percentage of the memory being used (paragraph 45). It would further be beneficial for a user to determine the maximum, such as taught by Gaddis, this allows the user to set further rules regarding route exporting, in addition to those as taught by Pesce.

11. As per claim 12, Sankaran further teaches:

an export limit indicative of a maximum number of routes that can be exported (paragraph 29, where a threshold may be established);

comparing the counted number of routes to the export limit (paragraph 35, where a threshold may be reached); and

rejecting additional routes exported when the counted number of routes exceeds the export limit (paragraph 45, where additional routes may be discarded once a threshold is reached).

12. As per claim 13, Pesce further teaches waiting for user intervention before accepting routes (column 7, lines 11-14, where a user may set rules governing route exportation).

13. As per claim 15, Sankaran further teaches updating routing information when the count exceeds the limit to clear the routes exported (paragraph 45, where routes may be deleted from the table when a limit is reached or exceeded).

14. As per claim 18, Pesce teaches a system comprising:

routes exported from an external routing protocol executing on a network device to an interior routing protocol executing on the network device in accordance with the export limit (column 2, lines 51-64, where the network element may contain multiple routing databases including those that may run both interior and exterior routing protocols, also column 3, lines 23-28, where the routing information may be exported from one database operating in one protocol to another database operating in another protocol); and

a plurality of instances of the interior routing protocol executing on the system (column 2, lines 51-64, where the network element may contain multiple routing databases including those that may run both interior and exterior routing protocols, also column 3, lines 23-28, where the routing information may be exported from one database operating in one protocol to another database operating in another protocol),

wherein the control unit separately exports to each of the instances (column 2, lines 51-64, where the network element may contain multiple routing databases including those that may run both interior and exterior routing protocols, also column 3, lines 23-28, where the routing information may be exported from one database operating in one protocol to another database operating in another protocol), and



wherein the control unit identifies an instance of the interior routing protocol to which routes were exported, accesses the respective rule, and rejects additional routes exported from the exterior routing protocol to the identified instance based on the comparison (column 7, lines 11-14, where a user may set rules governing route exportation).

Pesce does not expressly teach a client setting an export limit for the device. Sankaran teaches a method for managing discard algorithms comprising:

maintaining a count of routes exported (paragraph 35, where a threshold may be reached); and

rejecting additional routes exported when the count exceeds the export limit set by the command (paragraph 45, where additional routes may be discarded once a threshold is reached).

It would have been obvious to one of ordinary skill in the art at the time of the invention to utilize a route count and rejecting such as taught by Sankaran in a routing system such as taught by Pesce. Pesce's system generally allows for routes to be exported from one protocol operating on a network element to another protocol operating on a network element. Pesce further allows for blocking routes by a manager (column 4, lines 52-61). It would be beneficial in such a system to utilize a threshold and count of routes for blocking, as the routing manager may operate under blocking conditions after a threshold of a table is reached, such as a percentage of the memory being used (paragraph 45).

Neither Pesce nor Sankaran teaches a user defined size threshold. Gaddis teaches a method for improving databases comprising:

receiving a prefix limit command from a client to specify a limit for a table size (column 17, lines 29-34, where a prefix limit may be set for injected routes).

It would have been obvious to utilize a user defined limit such as that taught by Gaddis in a system for exporting routes such as that taught by either Pesce or Sankaran.

Pesce's system generally allows for routes to be exported according to definitions by a user, including those determining whether routes may be exported (column 7, lines 11-14). Pesce's system also allows a manager to determine when blocking conditions occur (column 4, lines 52-61). It would be beneficial to allow for a user to determine a maximum number of routes, such as taught by Sankaran, as the routing manager may operate under blocking conditions after a threshold of a table is reached, such as a percentage of the memory being used (paragraph 45). It would further be beneficial for a user to determine the maximum, such as taught by Gaddis, this allows the user to set further rules regarding route exporting, in addition to those as taught by Pesce.

15. As per claim 19, Gaddis further teaches a prefix counter to count the routes exported to the interior routing protocol and generate a prefix count, wherein the control unit compares the prefix count to the export limit and limits the number of routes exported to the interior routing protocol based on the comparison (column 17, lines 29-34, where a prefix limit may be set for injected routes).

16. As per claim 20, Gaddis further teaches rejecting additional routes to be exported to the interior routing protocol when the prefix count exceeds the export limit (column

17, lines 29-34, where a prefix limit may be set for injected routes, and routes may be filtered by prefix count).

17. As per claim 21, no reference expressly teaches the exterior routing protocol supports a larger number of routes than the interior routing protocol. However, Pesce teaches that BGP may be utilized with OSPF on a network element. BGP protocols maintain routes between ASes, therefore is theoretically capable of maintaining routing information to all systems on a network. OSPF is used within an AS, and therefore maintains routing information only with a selection of computers on the greater network

18. As per claim 22, Gaddis further teaches communicating with an internet service provider via the exterior routing protocol (column 1, lines 20-27, where common connections utilize ISPs).

19. As per claim 26, Pesce further teaches the system comprises a router (Figure 1).

20. As per claim 33, Pesce teaches a method comprising:

directing a network device to export routes from an exterior routing protocol executing on the network device to an interior routing protocol executing on the network device (column 2, lines 51-64, where the network element may contain multiple routing databases including those that may run both interior and exterior routing protocols, also column 3, lines 23-28, where the routing information may be exported from one database operating in one protocol to another database operating in another protocol);

receiving an export rule for routes that may be exported from the exterior routing protocol to a specific instance of the interior routing protocol (column 7, lines 11-14, where a user may set rules governing route exportation);

exporting routes from the exterior routing protocol to the specific instance of the interior routing protocol (column 2, lines 51-64, where the network element may contain multiple routing databases including those that may run both interior and exterior routing protocols, also column 3, lines 23-28, where the routing information may be exported from one database operating in one protocol to another database operating in another protocol);

blocking routes from the exterior routing protocol if the rule determines the route should not be exported (column 4, lines 52-61, where the updated mappings may be blocked by the manager).

Pesce does not expressly teach a client setting an export limit for the device. Sankaran teaches a method for managing discard algorithms comprising:

maintaining a count of routes exported (paragraph 35, where a threshold may be reached); and

rejecting additional routes exported when the count exceeds the export limit set by the command (paragraph 45, where additional routes may be discarded once a threshold is reached).

It would have been obvious to one of ordinary skill in the art at the time of the invention to utilize a route count and rejecting such as taught by Sankaran in a routing system such as taught by Pesce. Pesce's system generally allows for routes to be exported from one protocol operating on a network element to another protocol operating on a network element. Pesce further allows for blocking routes by a manager (column 4, lines 52-61). It would be beneficial in such a system to utilize a threshold and count of

routes for blocking, as the routing manager may operate under blocking conditions after a threshold of a table is reached, such as a percentage of the memory being used (paragraph 45).

Neither Pesce nor Sankaran teaches a user defined size threshold. Gaddis teaches a method for improving databases comprising:

receiving a prefix limit command from a client to specify a limit for a table size (column 17, lines 29-34, where a prefix limit may be set for injected routes). It would have been obvious to utilize a user defined limit such as that taught by Gaddis in a system for exporting routes such as that taught by either Pesce or Sankaran. Pesce's system generally allows for routes to be exported according to definitions by a user, including those determining whether routes may be exported (column 7, lines 11-14). Pesce's system also allows a manager to determine when blocking conditions occur (column 4, lines 52-61). It would be beneficial to allow for a user to determine a maximum number of routes, such as taught by Sankaran, as the routing manager may operate under blocking conditions after a threshold of a table is reached, such as a percentage of the memory being used (paragraph 45). It would further be beneficial for a user to determine the maximum, such as taught by Gaddis, this allows the user to set further rules regarding route exporting, in addition to those as taught by Pesce.

21. As per claim 39, Pesce further teaches the management interface receives the command from a remote client (column 7, lines 11-14, where a user may set rules governing route exportation).

22. As per claim 40, Pesce further teaches the remote client comprises one of a human user and an automated script (column 7, lines 11-14, where a user may set rules governing route exportation).

23. Claims 6 and 32 are rejected under 35 U.S.C. 103(a) as being unpatentable over US 7 328 278, Pesce et al, US 2003/0231587, Sankaran et al, and US 7 554 930, Gaddis et al as applied to claims above, and further in view of US 6 212 188, Rochberger et al.

24. As per claim 6, neither reference teaches a method to update the routing information in response to a state change of the device. Rochberger teaches a method of routing in a network when a node is in overload state comprising:

updating routing information to set an overload bit of a link state prefix associated with the routes when the count exceeds the export limit (Column 2, lines 8-14, where the state information is contained in PTSE messages. Along with column 5, lines 30-35, where the overloaded node sends a PTSE message to other nodes notifying itself as being overloaded, it is inherent that the PTSE contains the overload information, and is changed when the node goes into overload status); and  
advertising the updated routing information to a network device (Column 5, lines 30-35, where the overloaded node sends a message to other nodes notifying itself as being overloaded).

It would have been obvious to one of ordinary skill in the art at the time of the invention to include an overload notification such as that described by Rochberger in a network

system such as that taught by Sankaran or Pesce. Pesce's system would benefit, as the routing protocols could notify others when they are reached, relieving the processing required to process routes. Sankaran's system would benefit, as it could notify other devices when the address limit within a table has been reached, relieving the processing required for rejecting additional addresses sent. The notification method described by Rochberger can be used in any system, as it only describes how a node reacts to being in an overload state, and does not affect the performance of the node in normal functions. This would allow the notification method to be used in any system, including that taught by Sankaran or Pesce.

25. As per claim 32, neither reference teaches a method to update the routing information in response to a state change of the device. Rochberger teaches a method of routing in a network when a node is in overload state comprising:

updating routing information to set an overload bit of a link state prefix associated with the routes when the count exceeds the export limit (Column 2, lines 8-14, where the state information is contained in PTSE messages. Along with column 5, lines 30-35, where the overloaded node sends a PTSE message to other nodes notifying itself as being overloaded, it is inherent that the PTSE contains the overload information, and is changed when the node goes into overload status); and

advertising the updated routing information to a network device (Column 5, lines 30-35, where the overloaded node sends a message to other nodes notifying itself as being overloaded).

It would have been obvious to one of ordinary skill in the art at the time of the invention to include an overload notification such as that described by Rochberger in a network system such as that taught by Sankaran or Pesce. Pesce's system would benefit, as the routing protocols could notify others when they are reached, relieving the processing required to process routes. Sankaran's system would benefit, as it could notify other devices when the address limit within a table has been reached, relieving the processing required for rejecting additional addresses sent. The notification method described by Rochberger can be used in any system, as it only describes how a node reacts to being in an overload state, and does not affect the performance of the node in normal functions. This would allow the notification method to be used in any system, including that taught by Sankaran or Pesce.

### ***Conclusion***

26. **THIS ACTION IS MADE FINAL.** Applicant is reminded of the extension of time policy as set forth in 37 CFR 1.136(a).

A shortened statutory period for reply to this final action is set to expire THREE MONTHS from the mailing date of this action. In the event a first reply is filed within TWO MONTHS of the mailing date of this final action and the advisory action is not mailed until after the end of the THREE-MONTH shortened statutory period, then the shortened statutory period will expire on the date the advisory action is mailed, and any extension fee pursuant to 37 CFR 1.136(a) will be calculated from the mailing date of the advisory action. In no event, however, will the statutory period for reply expire later than SIX MONTHS from the mailing date of this final action.



Any inquiry concerning this communication or earlier communications from the examiner should be directed to THOMAS RICHARDSON whose telephone number is (571) 270-1191. The examiner can normally be reached on Monday through Thursday, 8am-5pm EST.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, William Vaughn can be reached on (571) 272-3922. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300.

Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see <http://pair-direct.uspto.gov>. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free). If you would like assistance from a USPTO Customer Service Representative or access to the automated information system, call 800-786-9199 (IN USA OR CANADA) or 571-272-1000.

TR  
/William C. Vaughn, Jr./  
Supervisory Patent Examiner, Art Unit 2444